
QUANTUM HOMOMORPHIC ENCRYPTION

Yfke Dulek



(joint work with Christian Schaffner and Florian Speelman)



Institute for Logic, Language
and Computation (ILLC)
University of Amsterdam



Research Center for
Quantum Software



Centrum
Wiskunde & Informatica

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1. HOMOMORPHIC ENCRYPTION
 2. PREVIOUS RESULTS
 3. NEW RESULT
-

QUANTUM HOMOMORPHIC ENCRYPTION



KEY GENERATION



public key



secret key



evaluation key

quantum state



ENCRYPTION
(secure)



+

$|x\rangle$

\mapsto



quantum data



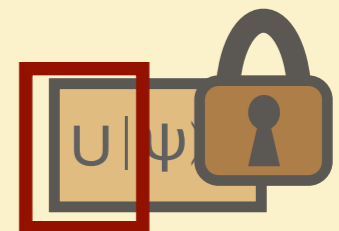
EVALUATION



+



\mapsto



quantum circuit



DECRYPTION
(compact)



+



\mapsto

$Uf(x)$

✓ HOMOMORPHIC ENCRYPTION

2. PREVIOUS RESULTS

3. NEW RESULT

PREVIOUS RESULTS: OVERVIEW

- Classical homomorphic encryption: solved [G09]
 - under (quantum-safe) computational assumptions (e.g. LWE)
- Quantum homomorphic encryption: only partial results

PREVIOUS RESULTS: OVERVIEW

	homomorphic for	compact?	security
Not encrypting	all circuits	yes	none
Quantum one-time pad	none	yes	inf theoretic
Append circuit description	all circuits	no: proportional to (# gates)	inf theoretic
Clifford Scheme	Clifford circuits	yes	computational
Clifford gates: [BJ15] AUX $P = [B, I, 5]: EPR$	circuits with constant T-depth	yes	$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$
$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$	all circuits	no: proportional to $(\# \text{ T-gates})^2$	
Also need: [OTF15]	circuit with constant # of T-gates	yes	inf theoretic
T = Our result $\begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix}$	all circuits of polynomial size (levelled QHE)	yes	computational

(comparison based on Stacey Jeffery's slides)

[BJ15] A. Broadbent, S. Jeffery. Quantum Homomorphic Encryption for Circuits of Low T-gate Complexity. CRYPTO 2015

[OTF15] Y. Ouyang, S-H. Tan, J. Fitzsimons. Quantum homomorphic encryption from quantum codes. [arxiv:1508.00938](https://arxiv.org/abs/1508.00938)

[YPDF14] L. Yu, C. Pérez-Delgado, J. Fitzsimons. Limitations on information-theoretically-secure quantum homomorphic encryption.



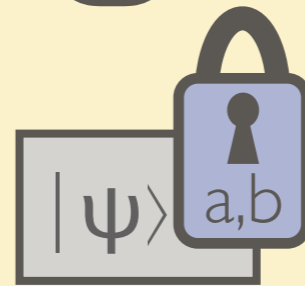
CLIFFORD SCHEME: P, H, CNOT

Ingredient 1: quantum one-time pad

encryption: pick $a, b \in_R \{0, 1\}$



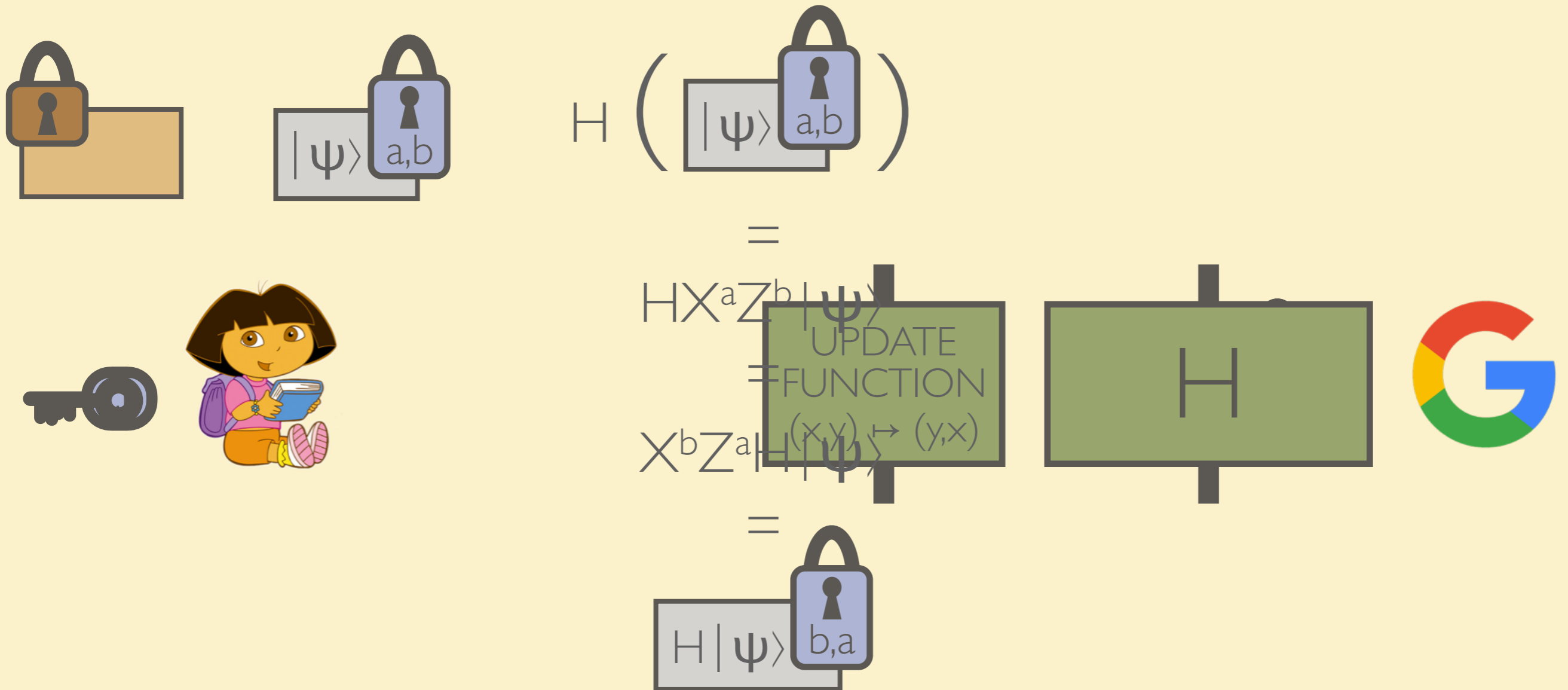
$$|\psi\rangle \mapsto X^a Z^b |\psi\rangle =$$



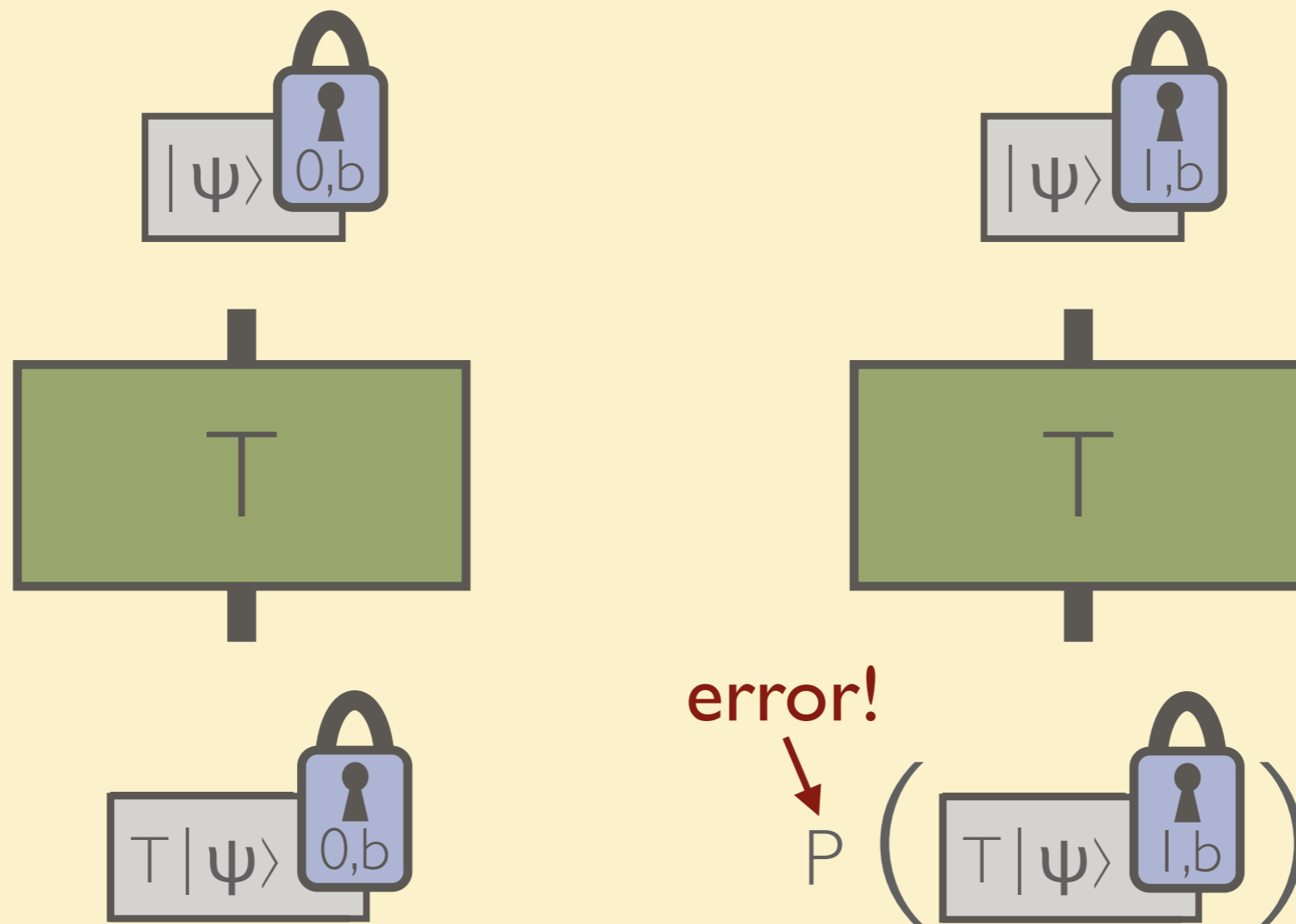
decryption: $X^a Z^b |\psi\rangle \mapsto |\psi\rangle$

Ingredient 2: classical homomorphic encryption

CLIFFORD SCHEME: P, H, CNOT



THE CHALLENGE: T GATE



how to apply correction P^{-1} iff $a = 1$?

✓ HOMOMORPHIC ENCRYPTION

✓ PREVIOUS RESULTS

3. NEW RESULT

NEW RESULT

	homomorphic for	compact?	security
Not encrypting	all circuits	yes	none
Quantum one-time pad	none	yes	inf theoretic
Append circuit description	all circuits	no: proportional to (# gates)	inf theoretic
Clifford Scheme	Clifford circuits	yes	computational
[BJ15]: AUX	circuits with constant T-depth	yes	computational
[BJ15]: EPR	all circuits	no: proportional to (# T-gates) ²	computational
[OTF15]	circuit with constant # of T-gates	yes	inf theoretic
Our result	circuits of polynomial size (levelled QFHE)	yes	computational

(comparison based on Stacey Jeffery's slides)

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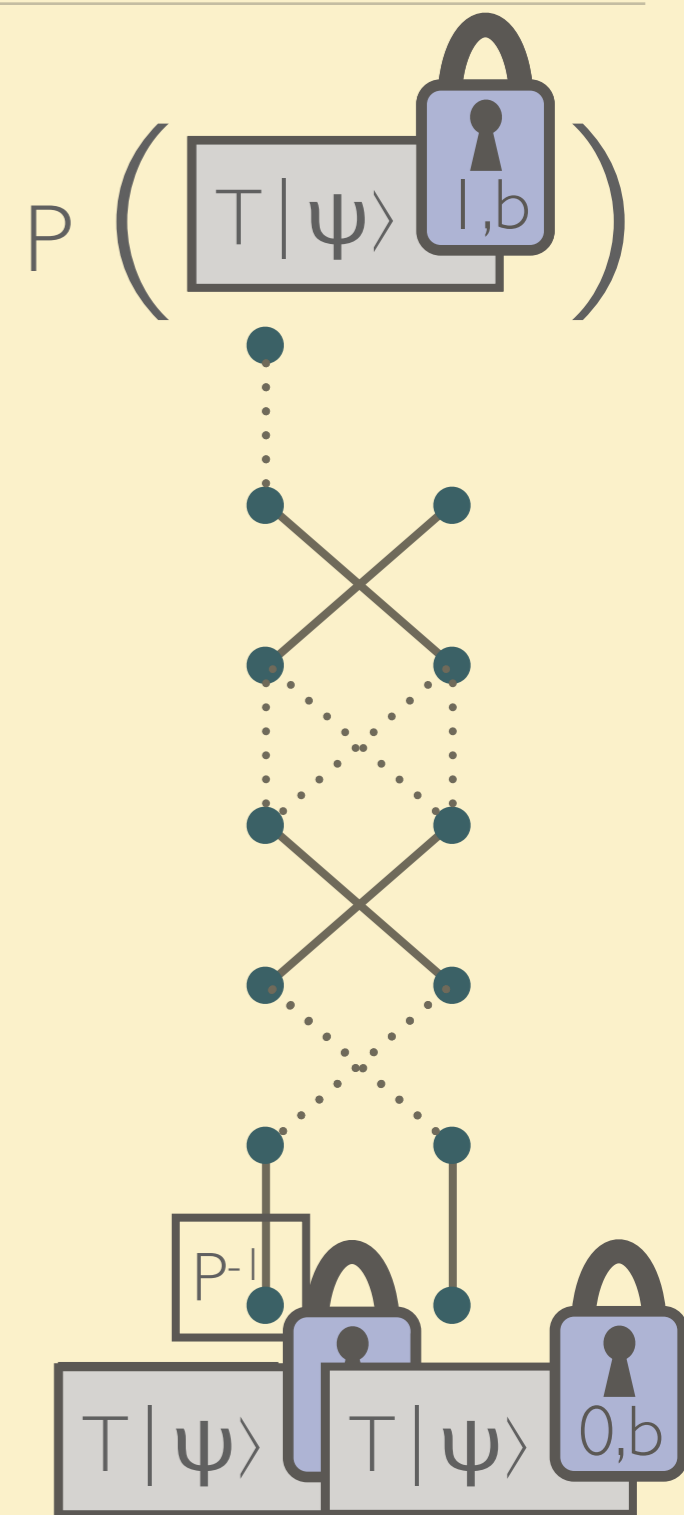


ERROR-CORRECTION GADGET

A quantum state (part of the evaluation key):

- Generation: entangle pairwise, according to 
- Usage: Bell measurements according to 

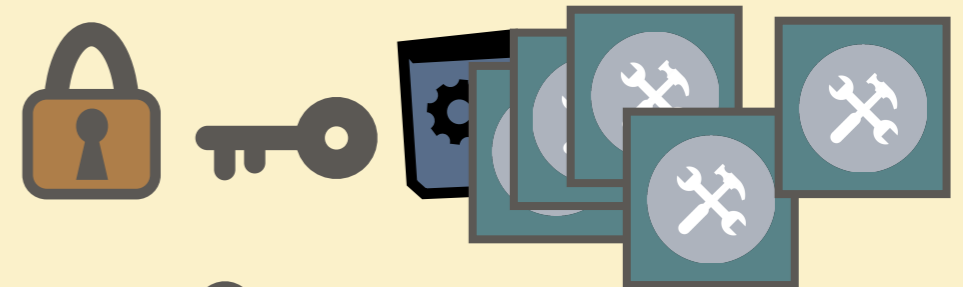
The gadget computes a permutation branching program for decrypt(, )



NEW SCHEME: OVERVIEW

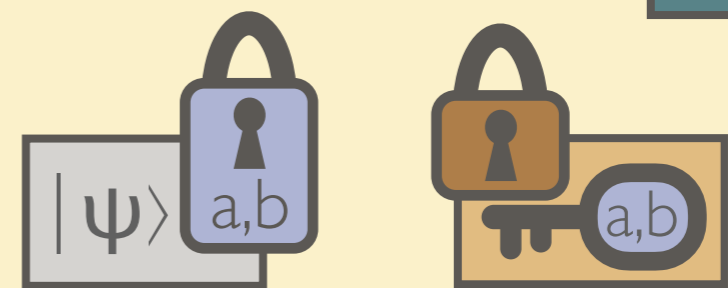
KEY GENERATION

- classical keys
- gadgets



ENCRYPTION

- apply quantum one-time pad
- classically encrypt pad keys

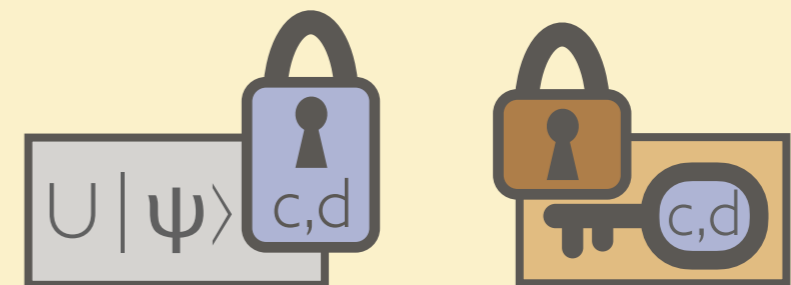


EVALUATION

- after **H** / **P** / **CNOT**: classically update keys
- after **T**: use 

DECRYPTION

- classically decrypt pad keys
- remove quantum one-time pad



FUTURE WORK

- non-leveled QFHE?
- multiparty quantum computation?
- quantum obfuscation?
- ...





THANK YOU!



QuSoft

