

Towards new upper bounds for $\text{QMA}(k)$

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What is known so far...

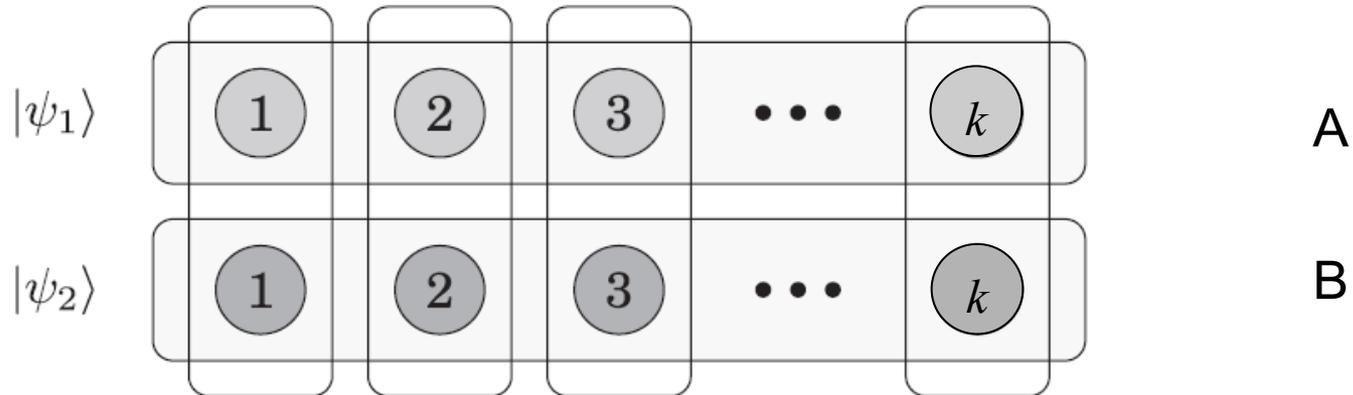
$$\text{QMA} \subseteq \text{QMA}(2) \subseteq \text{NEXP}$$

- only trivial bounds known so far for the general case
- under some assumptions, non-trivial bounds are known:
 - [BCY11], [SW15] show quasi-polynomial time algorithms (in d) for h_{sep} assuming $\|\cdot\|_F$ -norm bounds on measurement operators
 - both based on geometric methods (\mathcal{E} -nets, SDP hierarchies)
 - don't exploit circuit structure of verifier at all \rightarrow new angle?

Chailloux and Sattath's result

- Search for a QMA(2)-complete problem [CS11]
- Obvious candidate: **Local Hamiltonian (LH)**, but with *separable witnesses state*, called **Separable Local Hamiltonian (SLH)**
- Turns out, **Separable Local Hamiltonian** is in fact QMA-complete!
- Nevertheless, [CS11] show **Separable Sparse Hamiltonian (SSH)** is to be QMA(2)-complete
- **Sparse vs. Local** - where is the big difference?

Harrow-Montanaro normal form



- every $\text{QMA}(k)$ verifier can be reduced to a $\text{QMA}(2)$ -verifier [HM13], such that:
 - the two witness states in the A:B subsystems are identical copies
 - only k -many $\text{poly}(n)$ -qubit SWAP tests are performed across A:B
 - the original verifier circuit is only executed on either A or B
 - in the YES case, measurement is *separable* operator across A:B!

Separable Sparse Hamiltonian

$$H = H_{in} + H_{prop} + H_{out} \quad H_{prop} = \sum_{t=1}^T H_t$$

$$H_t = -\frac{1}{2}|t\rangle\langle t-1| \otimes U_t - \frac{1}{2}|t-1\rangle\langle t| \otimes U_t^\dagger + \frac{1}{2}(|t\rangle\langle t| + |t-1\rangle\langle t-1|) \otimes I.$$

- [HM13] verifier circuit mapped to a Local Hamiltonian yields:
 - k groups of $poly(n)$ -many local terms encoding CSWAP unitaries across the A:B boundary at the beginning of the computation
 - $poly(n)$ -many terms acting only on one subsystem, encoding the verifier
- Problem: history state of this Hamiltonian is NOT separable!
 - single-qubit CSWAPs at different times create entanglement
- Solution: CSWAP all qubits of each group in a “single time step“!
 - witness state now separable, but k terms now sparse, but non-local

My starting point

$$H_{\text{CSWAP}} = -\frac{1}{2} \overbrace{\text{SWAP}_{1,\ell+1} \otimes \text{SWAP}_{2,\ell+2} \otimes \cdots \otimes \text{SWAP}_{\ell,2\ell}}^{\text{swap terms } S_i} \otimes \overbrace{|1\rangle\langle 1|}^{\text{control } C} \otimes \overbrace{(|t\rangle\langle t-1| + |t-1\rangle\langle t|)}^{\text{time propagation } T}$$

- Can we simulate such terms (more) locally?
- Is there a “parallel Feynman-Kitaev construction“ that “executes“ parallel unitaries in the same “time step“ in the history state?
- Idea: use perturbation theory gadgets to reduce locality!
 - used before to reduce 5-local to 2-local QMA-complete Hamiltonians
 - exponential cost in Hamiltonian norm
- A reduction in promise gap might be acceptable, though...

Perturbation theory review

- Start with a simple, unperturbed Hamiltonian with a large spectral gap Δ between ground state and excited states, acting on ancilla w :

$$H = \Delta |1\rangle\langle 1|_w$$

- H acts trivially on other qubits \rightarrow degenerate GS
- Adding a perturbation V with $\|V\| < \Delta/2$, can (under certain conditions) approximate the spectrum of a desired effective Hamiltonian H_{eff} in the low-energy spectrum ($E < \Delta/2$) of

$$\tilde{H} = H + V$$

Perturbation theory review

Theorem 10 (Eigenvalue-Approximating Perturbation Theorem [OT08, Theorem 7]). *Let $\|V\| \leq \Delta/2$ where Δ is the spectral gap of H and $\lambda(H) = 0$. Let $\tilde{H}|_{<\Delta/2}$ be the restriction of $\tilde{H} = H + V$ to the space of eigenstates with eigenvalues less than $\Delta/2$. Let there be an effective Hamiltonian H_{eff} with $\text{Spec}(H_{\text{eff}}) \subseteq [a, b]$. If the self-energy $\Sigma_-(z)$ for all $z \in [a - \varepsilon, b + \varepsilon]$ where $a < b < \Delta/2 - \varepsilon$ for some $\varepsilon > 0$, has the property that*

$$\|\Sigma_-(z) - H_{\text{eff}}\| \leq \varepsilon, \quad (5)$$

then each eigenvalue $\tilde{\lambda}_j$ of $\tilde{H}|_{<\Delta/2}$ is ε -close to the j th eigenvalue of H_{eff} . In particular

$$|\lambda(H_{\text{eff}}) - \lambda(\tilde{H})| \leq \varepsilon. \quad (6)$$

- Generic perturbation theorem for any choice of V and H_{eff} , where the
- *Self-energy expansion* $\Sigma_-(z)$ of \tilde{H} converges to H_{eff}
- If gadget V is more local than H_{eff} \rightarrow *effective locality reduction!* [OT08]
- Alternatives: *Bloch expansion* or *Schrieffer-Wolff transformation*
- Observation: No statement at all about eigenvectors!

Perturbation theory review

Theorem 11 (Norm-Approximating Perturbation Theorem [OT08, Theorem A.1]). *Given is a Hamiltonian H such that no eigenvalues of H lie between $\lambda_- = \lambda_* - \Delta/2$ and $\lambda_+ = \lambda_* + \Delta/2$. Let $\tilde{H} = H + V$ where $\|V\| \leq \Delta/2$. Let there be an effective Hamiltonian H_{eff} with $\text{Spec}(H_{\text{eff}}) \subseteq [a, b]$, $a < b$. We assume that $H_{\text{eff}} = \Pi_- H_{\text{eff}} \Pi_-$. Let D_r be a disk of radius r in the complex plane centered around $z_0 = \frac{b+a}{2}$. Let r be such that $b + \varepsilon < z_0 + r < \lambda_*$. Let $w_{\text{eff}} = \frac{b-a}{2}$. Assume that for all $z \in D_r$ we have $\|\Sigma_-(z) - H_{\text{eff}}\| \leq \varepsilon$. Then*

$$\|\tilde{H}_{<\lambda_*} - H_{\text{eff}}\| \leq \frac{3(\|H_{\text{eff}}\| + \varepsilon)\|V\|}{\lambda_+ - \|H_{\text{eff}}\| - \varepsilon} + \frac{r(r + z_0)\varepsilon}{(r - w_{\text{eff}})(r - w_{\text{eff}} - \varepsilon)}. \quad (7)$$

- But even more is true: approximation in *operator norm!*
- Implies that original witness space is close to perturbed one
- Original 5-LH to 2-LH reductions only cared about eigenvalues
 - structure of eigenvectors of no relevance – only promise gap important

Subdivision Gadget

Concrete choice of V to reduce locality, also works in parallel [OT08]:

$$H = \Delta |1\rangle\langle 1|_w, \quad V = H'_{\text{else}} + \sqrt{\Delta/2} (-A + B) \otimes X_w$$
$$\tilde{H} = H + V$$

(l+1)-local

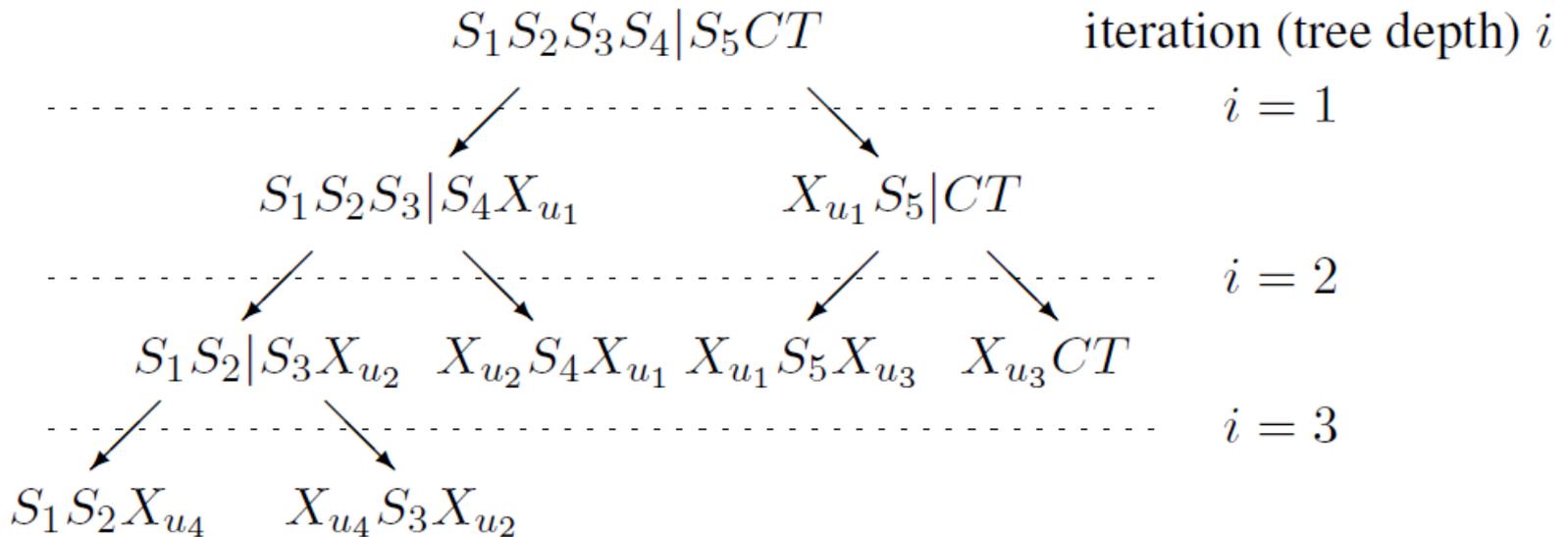
$$H_{\text{target}} = H'_{\text{else}} + A \otimes B \leftarrow 2l\text{-local}$$

$$H_{\text{eff}} = H_{\text{target}} \otimes |0..0\rangle\langle 0..0|$$

Lemma: If $\Delta > \text{poly}(n, k) / \varepsilon^2$, then $\left\| \tilde{H}|_{<\Delta/2} - H_{\text{target}} \otimes |0..0\rangle\langle 0..0| \right\| \leq \varepsilon$

Subdivision Gadget

- Breaking l -local down to 3-local terms [CBBK15]:



- Exponential scaling of gap Δ in l [BDLT08]:

$$\Delta_{i+1} = \frac{\text{poly}(n, k)}{\varepsilon^2} \Delta_i^3 \quad \Delta = \left(\frac{\text{poly}(n, k)}{\varepsilon^2} \right)^{3^{\log(\ell)}} = \left(\frac{\text{poly}(n, k)}{\varepsilon^2} \right)^{\text{poly}(\ell)} = 2^{O(\text{poly}(\ell) \log(nk/\varepsilon))}$$

QMA(2): **SSH** to **SLH** reduction

- Applying the subdivision gadget to our **SSH** instance, we can break down the k non-local but sparse terms to 3-local ones, yielding **SLH**
 - Definition of **Local Hamiltonian** requires poly-bounded norm of H , but this reduction increases it exponentially
- Solution: rescale **SLH** instance dividing by Δ
- Operator norm of H now $O(1)$
 - Soundness/completeness divided by $\Delta \rightarrow$ *promise gap* divided by Δ
- SSH to SLH conversion at the cost of exp. small promise gap! (...)

Deciding SLH using ε -nets

Corollary 6 ([SW15, Problem 3, Corollary 6]). Take the expression $Q = \sum_{i=1}^r H_i$ of any l -local Hamiltonian over $A_1 \otimes \cdots \otimes A_k$ (each A_i is of dimension $d = 2^n$) such that $\|H_i\|_{\text{op}} \leq w$ for each i as input. Assuming $k, l = O(1)$, the quantity

$$\text{OptSep}(Q) = \min \langle Q, \rho \rangle \text{ subject to } \rho \in \text{SepD}(A_1 \otimes \cdots \otimes A_k) \quad (1)$$

where $\text{SepD}(A_1 \otimes \cdots \otimes A_k)$ is the set of separable density operators over the space $A_1 \otimes \cdots \otimes A_k$, can be approximated to precision δ in

$$\text{DTIME}(\exp(O(\log^{O(1)}(d)(\log \log(d) + \log(w/\delta)))) \times \text{poly}(d, w, 1/\delta)), \quad (2)$$

which is quasi-polynomial in $d, w, 1/\delta$. If n is considered as the input size and $w/\delta = O(\text{poly}(n))$, then $\text{OptSep}(Q)$ can be approximated to precision δ in PSPACE.

→ yields an **EXP** algorithm for $d=2^n$ and $\delta=1/2^{\text{poly}(n)}$, thus “QMA(2) \subseteq EXP”

→ if there weren't a problem lurking...

- **The Problem:** the perturbation theory requires **projection onto the low-energy** subspace!

$$\left\| \tilde{H}_{<\Delta/2} - H_{\text{target}} \otimes |0\dots 0\rangle\langle 0\dots 0| \right\| \leq \varepsilon$$

- Projected Hamiltonian: no longer local (self-defeating perturbation)
- Projected separable states: close in energy, but no longer separable
- Unprojected separable states: **overlap with high-energy subspace**
- **Soundness OK**: high energy states stay high energy
- **Completeness NOK**: original witness has too much energy

Analyzing the issue

Geometric picture: Schrieffer-Wolff transformation [BDL11, BLDT08]

$$H_0 = \Delta|1\rangle\langle 1|_u, \quad V = \sqrt{\Delta J/2} X_u \otimes (-A + B) + V_{\text{extra}} \quad S = -i\sqrt{\frac{J}{2\Delta}} Y_u \otimes (-A + B)$$

$$e^S H e^{-S} = \left(H + [S, H] + \frac{1}{2}[S, [S, H]] \right) + O(J^{3/2}\Delta^{-1/2}) = \begin{bmatrix} H_{\text{target}} & 0 \\ 0 & \Delta I + O(J) \end{bmatrix} + O(J^{3/2}\Delta^{-1/2}).$$

→ H_{target} is reproduced in a slightly rotated basis (up to some error)

→ overlap with excited subspace is penalized by Δ

How much overlap with excited space $(1-P)$ of H can an ground state of H_0 have?

Lemma: Let P_0 be projector on eigenvectors of H_0 . Let P be projector on

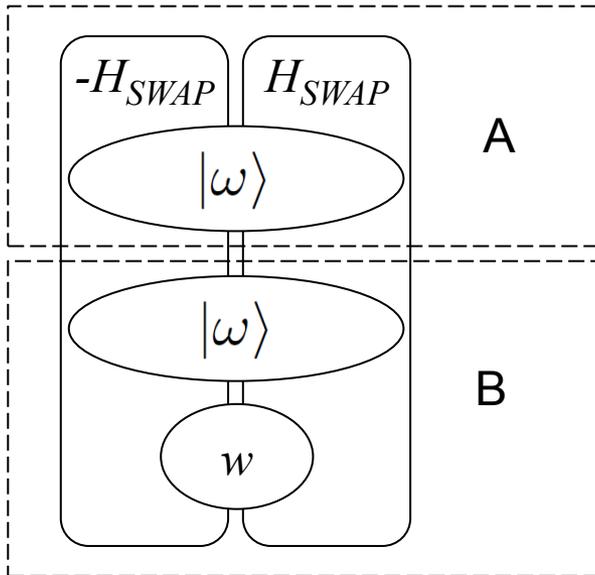
$(<\Delta/2)$ -eigenvectors of H . Then, $\|P - P_0\| \leq \frac{2J\|V\|}{\Delta} = \sqrt{\frac{2J^3}{\Delta}}$

$O\left(\frac{1}{\sqrt{\Delta}}\right)$ overlap with $(1 - P)$

→ Thus, upper bound on energy: $\langle \psi | \langle 0 | H | \psi \rangle | 0 \rangle \leq O(\sqrt{\Delta})$

$O(\Delta)$ energy penalty on $(1 - P)$

Example for violation



$$\tilde{H} = H + V \quad H = \Delta|1\rangle\langle 1|_w \quad \Delta = 100$$

$$V = \mathbb{1} - \underbrace{\sqrt{\frac{\Delta}{2}} \text{SWAP}_1 \otimes X_w}_{H_{SWAP}} + \underbrace{\sqrt{\frac{\Delta}{2}} \text{SWAP}_2 \otimes X_w}_{H_{SWAP}}$$

$$H_{\text{eff}} = -\text{SWAP}_1 \otimes \text{SWAP}_2 \otimes |0\rangle\langle 0|_w$$

$$|\psi\rangle = |\omega\rangle \otimes |\omega\rangle \otimes |0\rangle \quad |\omega\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

- $P|\psi\rangle$ necessarily has small superposition on ancilla
- superposition on ancilla+B doesn't really help much

$$\langle \psi | H_{\text{eff}} | \psi \rangle = -1$$

$$\langle \psi | P \tilde{H} P | \psi \rangle = -0.94$$

$$\langle \psi | \tilde{H} | \psi \rangle = +1$$

- Find new “parallel” Feynman-Kitaev Hamiltonians, e.g.
 - implement one time step with k ancillas as parallel clocks
 - each clock ancilla couples to one of the CSWAP gates (\rightarrow local)
 - couple k ancillas with n.n. $-ZZ$ terms to enforce synchronized tick
 - add small local $-X$ terms to single out $|+\rangle = |00\dots 0\rangle + |11\dots 1\rangle$
 - \rightarrow Ising model ($ZZ+X$ terms): has analytical solution!
 - similar to Jordan-Farhi [JF08] direct k -local to 2-local gadget
- $\rightarrow -X$ needed to single out $|+\rangle = |00\dots 0\rangle + |11\dots 1\rangle$ as ground state
- \rightarrow otherwise, degenerate ground space, includes no time propagation
- \rightarrow but X acts like perturbation: non-zero amplitudes in GS, i.e. partially flipped clocks (non-separable), original witness energy too high

QMA(2) in PSPACE?

- Assume the reduction to $\mathbf{SLH}_{\text{gap}}$ worked, thus $\text{QMA}(2) \subseteq \text{EXP}$
- [CS11] also show that \mathbf{SLH} (poly-gap) is QMA-complete
 - Reduces to *Consistency of Density Matrices* (CDM) problem
 - Witness consists of classical density matrices and quantum consistency proof
 - CDM uses QMA+ verifier which requires $O(\text{poly}(1/\varepsilon))$ copies of witness state
- If we could show that $\mathbf{SLH}_{\text{gap}}$ is in QMA_{exp} , then we can use $\text{QMA}_{\text{exp}} = \text{PSPACE}$ [FL16] to put $\text{QMA}(2) \subseteq \text{PSPACE}$
- What about a very direct reduction of $\text{QMA}(2)$ to QMA_{exp} ?
 - E.g. receive one copy and “guess” the second one with exp. small chance?

Conclusion

SSH to **SLH** reduction (with small gap):

- ✗ Subdivision gadget seems to be self-defeating for $\text{QMA}(2)$
- Does not rule out alternative gadget designs, maybe very specific one
- Does not rule out alternative “parallel” Feynman-Kitaev constructions, that maintain locality by construction

Solving **SLH** with small gap:

- ✓ in EXP using ε -net algorithm of Shi and Wu
- might even be in PSPACE by reduction to $\text{QMA}_{\text{exp}} = \text{PSPACE}$

Thank you!